Efficient decoding for the Hayden-Preskill protocol

(Joint work with Beni Yoshida)

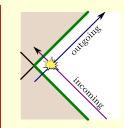
arXiv:1710.03363

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Motivation: quantum black holes

• Quantum fields in classical space (Hawking): $T = \frac{\varkappa}{2\pi}$, $S = \frac{A}{4}$

- Information scrambling: $\tau_{\rm scr} \approx (2\pi T)^{-1} \ln S$
 - Gravitational interaction between incoming and outgoing radiation
 - Dray-t'Hooft shock waves
 - OTOCs: $\langle W(t) Y(0) Z(t) X(0) \rangle$



- Evolution over Page's time, when half of the black hole evaporates
 - Full quantum gravity



Assumptions

• Thermal state is replaced with the maximally mixed state on a "typical subspace" \mathcal{L} :

$$\rho = Z^{-1}e^{-H/T} \longrightarrow \rho = \frac{I_{\mathcal{L}}}{d}, \quad d = \dim \mathcal{L} = \mathscr{E}^{\mathcal{S}} 2^{S}$$

• Late-time OTOCs (= almost perfect scrambling):

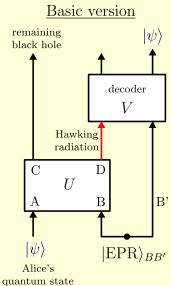
$$\langle W(t) Y(0) Z(t) X(0) \rangle$$

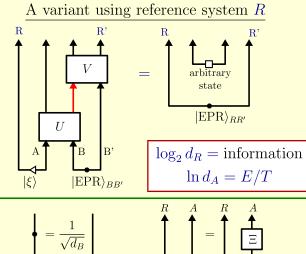
$$\approx \langle WZ \rangle \langle Y \rangle \langle X \rangle + \langle W \rangle \langle Z \rangle \langle YX \rangle - \langle W \rangle \langle Z \rangle \langle Y \rangle \langle X \rangle$$

where
$$W(t) = U^{\dagger}WU$$
, $Z(t) = U^{\dagger}ZU$, $Y(0) = Y$, $X(0) = X$

- Holds for a Haar-random unitary U
- Broadly applicable if X, Y, Z, W act on small subsystems

The Hayden-Preskill problem





 $(\Xi: R' \to A)$

Tensor diagrams

- Nodes are tensors; lines are index contractions.
- Time goes up.
 - Vertical sections of lines are associated with Hilbert spaces. If a line goes up and then down, the Hilbert space changes to the dual space.
 - Let $\psi \in A$ be a vector with elements c_i and $\psi^* \in A'$ a vector with elements c_i^* . Then

$$|\psi\rangle = \prod_{A}^{A} \psi \qquad |\psi^*\rangle = \prod_{A}^{A} \psi^* \qquad \langle\psi| = \prod_{A}^{\psi^*} \psi^* \qquad \langle\psi^*| = \prod_{A}^{\psi} \psi^* \psi^* = \prod_{A}^{\psi} \psi^* = \prod_{A}^{\psi} \psi^* \psi^* = \prod_{A}^{\psi} \psi^* =$$

$$-X^T$$
 is X upside-down:
$$\begin{bmatrix} A \\ X \end{bmatrix} = \begin{bmatrix} A \\ X^T \end{bmatrix} = \sum_{j,k} X_{jk} |j,k\rangle$$

When is the decoding possible?

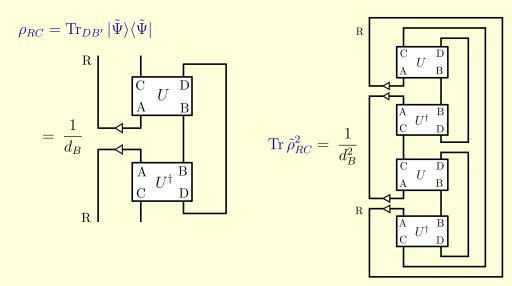
State of the world:
$$|\widetilde{\Psi}\rangle = \bigcup_{|\xi\rangle} U$$
 $\widetilde{\rho} = |\widetilde{\Psi}\rangle\langle\widetilde{\Psi}|$

- Black hole has "forgotten" Alice's secret $\Leftrightarrow \tilde{\rho}_{RC} \approx \tilde{\rho}_R \otimes \tilde{\rho}_C$
- Quantitative condition: Let

$$\delta = d_R d_C \operatorname{Tr} \tilde{\rho}_{RC}^2 - 1 \qquad (\delta \geqslant 0).$$

If $\delta \ll 1$, then Alice's secret can, in principle, be recovered from the Hawking radiation D and the purifying subsystem B'. In our algorithms (and in the original Hayden-Preskill work), δ determines the decoding fidelity.

Calculation of $1 + \delta = d_R d_C \operatorname{Tr} \tilde{\rho}_{RC}^2$



$= \mathrm{OTOC}(L, M)'$ $\delta = d_A d_R \Delta - 1, \qquad \Delta =$

Expression for the fidelity parameter δ

Generalizes the result of P. Hosur, X.-L. Qi, D. Roberts, B. Yoshida, arXiv:1511.04021)
$$L = d_A \bigcap_{i=1}^{A'} \bigcap_{j=1}^{A} \sum_{j=1}^{A} Y_j^T \otimes X_j, \qquad M = \bigcap_{j=1}^{D} \bigcap_{j=1}^{D'} \sum_{k=1}^{D} W_k \otimes Z_k^T$$

(Generalizes the result of P. Hosur, X.-L. Qi, D. Roberts, B. Yoshida, arXiv:1511.04021)
$$L = d_A \bigwedge_{A'}^{A'} \bigwedge_{A}^{A} = \sum_j Y_j^T \otimes X_j, \qquad M = \bigcup_{D=D'}^{D} \sum_{D'}^{D'} = \sum_k W_k \otimes Z_k^T$$

 $OTOC(L, M) = \sum_{j,k} \frac{1}{d} Tr((U^{\dagger}W_k U) Y_j (U^{\dagger}Z_k U) X_j)$

The late-time case Assumption:

 $\overline{\langle W(t) Y(0) Z(t) X(0) \rangle} \approx \langle WZ \rangle \langle Y \rangle \langle X \rangle + \langle W \rangle \langle Z \rangle \langle YX \rangle - \langle W \rangle \langle Z \rangle \langle Y \rangle \langle X \rangle$

$$UWU^{\dagger}$$
 UZU^{\dagger}

$$\underline{\text{Used in calculations}}: \quad \langle Y \rangle \langle X \rangle = \boxed{Y^T} \stackrel{X}{X} \qquad \langle YX \rangle = \boxed{Y^T} \stackrel{X}{X}$$

Result:
$$\Delta \approx \frac{1}{d_A d_R} + \frac{1}{d_D^2} - \frac{1}{d_A d_R d_D^2} \Rightarrow \delta = d_A d_R \Delta - 1 \leqslant \frac{d_A d_R}{d_D^2}$$

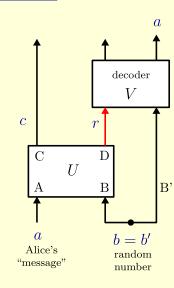
How hard is the decoding?

• The complexity is at least linear in d_R (i.e. exponential in the message size). Indeed, let $d_A = d_R$ and consider the classical case:

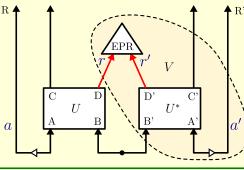
$$(c,r) = u(a,b)$$
discarded known to Bob

Thus, r = f(a), where f is random. The only general way to reconstruct a from r is exhaustive search.

• We show that the complexity is $O(d_A d_R C)$, where C is the size of the circuit for U.



Probabilistic decoder



Classically, it's just random guessing. The Hawking radiation is r = f(a); Bob picks a random a', computes r' = f(a'), and

compares it with r.

succeeds with probability
$$\sqrt{|\mathbf{J}_{1}|} (I_{-1} \otimes P_{-1}) \mathbf{J}_{1} = \Delta > (d_{-1}d_{-1})^{-1}$$

 $\langle \Psi_{\rm in}(I_{RC} \otimes P_D) \Psi_{\rm in} \rangle = \Delta \geqslant (d_A d_R)^{-1}$

Fidelity of probabilistic decoding

 $|\Psi_{\rm out}\rangle = \frac{1}{\sqrt{\Delta}} (I_{RC} \otimes P_D) |\Psi_{\rm in}\rangle = \frac{1}{\sqrt{\Delta}} |U|^{-1} U^*$ Projected state:

Fidelity:

$$F = /\Psi$$

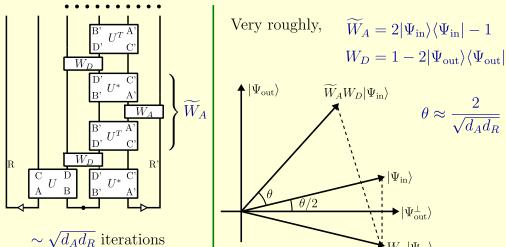
$$F = \langle \Psi_{\text{out}} | P_R | \Psi_{\text{out}} \rangle = \Delta^{-1} \langle \Psi_{\text{in}} | \underbrace{P_R (I_{RC} \otimes P_D)}_{\text{In}} | \Psi_{\text{in}} \rangle \geqslant \frac{1}{d_A d_R \Delta} = \boxed{\frac{1}{1 + \delta}}$$

$$P_R(I_{RC} \otimes P_D)$$
 | $P_R(I_{RC} \otimes P_D)$ |

$$|\Psi_{
m in}
angle\geqslantrac{1}{d_Ad_R\Delta}=$$

Deterministic decoder

• Uses Grover search to turn $|\Psi_{\rm in}\rangle$ to $|\Psi_{\rm out}\rangle$ without projection



More accurate description of the algorithm

1) Apply
$$U^*$$
 to produce $|\Psi_{\rm in}\rangle=$

2) Let
$$P_{D} = \bigcap_{D = D' = C' = R'}^{D = D' = C' = R'} P_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = B' = A' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'} \widetilde{P}_{A} = \bigcap_{D = B' = A' = R'}^{D = D' = C' = R'}$$

$$W_D = 1 - 2P_D,$$
 $\widetilde{W}_A = 2\widetilde{P}_A - 1$

$$W_D = 1 - 2P_D,$$
 $\widetilde{W}_A = 2\widetilde{P}_A - 1$
3) Apply $\widetilde{W}_A W_D$ repeatedly $\frac{\pi}{2\theta_*}$ times, where $\theta_* = 2\arcsin\left((d_A d_R)^{-1/2}\right)$.

• Let $\widetilde{P}_A P_D \widetilde{P}_A = \Pi = \sum_{i=1}^r \alpha_j |\psi_j\rangle \langle \psi_j|$, $|\Psi_{\rm in}\rangle = \sum_{j=1}^r \sqrt{p_j} \underbrace{|\eta_j\rangle_{RC} \otimes |\psi_j\rangle}_{|\Psi_j\rangle}$

eigenvalue decomposition, $\alpha_i > 0$

Then each vector $|\Psi_i\rangle$ evolves under $I_{RC}\otimes (W_AW_D)^m$ in a two-dimensional subspace with basis vectors $|\Phi_i\rangle$, $|\Phi_i^{\perp}\rangle$.

Analysis of the algorithm

$$|\Psi(m)\rangle = \sum_{j=1}^{r} \sqrt{p_j} \left(\sin\left(\left(m + \frac{1}{2}\right)\theta_j\right) |\Phi_j\rangle + \cos\left(\left(m + \frac{1}{2}\right)\theta_j\right) |\Phi_j^{\perp}\rangle \right)$$
where $\theta_j = 2 \arcsin\sqrt{\alpha_j}$

• We show that
$$r \leqslant d_R d_C$$
, $\sum_{i=1}^r \alpha_i = \frac{d_C}{d_A}$, $\sum_{i=1}^r \alpha_i^2 = \frac{d_C}{d_A} \Delta$.

 $r \leqslant d_R d_C, \quad \sum_{j=1}^r \alpha_j = \frac{d_C}{d_A}, \quad \sum_{j=1}^r \alpha_j^2 = \frac{d_C}{d_A} \Delta.$

If $\delta = d_A d_R \Delta - 1 = 0$ (ideal case), then $\alpha_j = (d_A d_R)^{-1/2}$ for all j.

Analysis of the algorithm (cont.)

• Let $\delta = d_A d_R \Delta - 1$, $m_* = \pi/(2\theta_*)$, where $\theta_* = 2\arcsin((d_A d_R)^{-1/2})$.

Then
$$(m_* + \frac{1}{2})\theta_j \approx \frac{\pi}{2}$$
,

$$|\Psi(m_*)\rangle \approx \sum_{j=1}^r \sqrt{p_j} |\Phi_j\rangle \approx |\Psi_{\text{out}}\rangle;$$
 the Euclidean distance is $O(\sqrt{\delta})$.

• <u>Conclusions</u>:

- The algorithm involves $O(\sqrt{d_A d_R})$ applications of U^* and U^T .
- The fidelity of the reconstructed state $|\Psi(m_*)\rangle$ is $O(\delta)$. Recall that in the case of almost perfect scrambling,

$$\delta \leqslant \frac{d_R d_A}{d_D^2}$$

Open questions

1. How to generalize the algorithm to thermal density matrices? We can do it under these unrealistic assumptions:

$$\rho_{AB} = \rho_A \otimes \rho_B, \qquad \rho_{CD} = \rho_C \otimes \rho_D, \qquad \rho_{CD} = U \rho_{AB} U^{\dagger}.$$

- 2. In the traversible wormhole story, (Gao, Jafferis, Wall 2016; Maldacena, Stanford, Yang 2017), the decoding happens in one go. What are the necessary/sufficient conditions in terms of OTOCs?
- 3. The Grover iterations bear some similarity with multiple shocks (Shenker, Stanford 2014). What is the exact relation?